# Principles of Programming Languages Lecture 12 

Wael Aboulsaadat wael@cs.toronto.edu

http://portal.utoronto.ca/

Acknowledgment: parts of these slides are based on material by Diane Horton \& Eric Joanis @ UoT References: Scheme by Dybvig

PL Concepts and Constructs by Sethi
Concepts of PL by Sebesta ML for the Working Prog. By Paulson
Prog. in Prolog by Clocksin and Mellish PL Pragmatics by Scott

## ML: introduction

- Developed at Edinburgh (early '80s) as Meta-Language for a program verification system
- Now a general purpose language
- There are two basic dialects of ML
- Standard ML (1991) \& ML 2000
- Caml (including Objective Caml, or OCaml)
- A pure functional language
- Based on typed lambda calculus
- Grew out of frustration with Scheme!
- Serious programs can be written without using variables

```
1930's Lambda calculus (Church)
1950's Lisp (McCarthy)
1960's semantics, deconstruction
1970's FP (Backus)
1980's Miranda (Turner), ML (Milner)
1990's Haskell
```

- Widely accepted
- reasonable performance (claimed)
- can be compiled
- syntax not as arcane as Scheme


## ML: main features

- Strong, static typing
- Quite a fancy type system!
- Parametric polymorphism
- Similar to OOP (in fact, it influenced OO)
- Pattern matching
- Function as a template
- Exception handling
- Allow you to handle errors/exception
- Type inference
- Recursive data type


## ML: how far have PL advanced?

- Writing a gcd implementation

$$
\operatorname{gcd}(m, n)=\left\{\begin{array}{cc}
n & m=0 \\
\operatorname{gcd}(n \bmod m, m) & m>0
\end{array}\right.
$$

Pascal

```
function gcd(m,n: integer): integer;
var prevm: integer;
begin
        while m<>0 do begin
            prevm := m; m := n mod m; n := prevm
        end;
        gcd := n
end;
```

Scheme
(define gcd
(lambda (m n)
(if (zero? n ) m
$(\operatorname{gcd} \mathrm{n}($ remainder m n$))$ )) $)$

ML
fun $\operatorname{gcd}(m, n)=$
if $\mathrm{m}=0$ then n else gcd(n mod m, m);

## ML: types \& expressions

- Primitive types
- bool, int, real, string
- Complex types
- list, tuple, array, record, function
- Each ML expression has a type associated with it.
- Interpreter builds the type expression for each input
- Cannot mix types in expressions

$$
2+3.0 \rightarrow \text { error! }
$$

- Must explicitly coerce/type-case e.g.

$$
\text { real }(2)+3.0 \text { : real }
$$

## ML: Primitive Types

- int
e.g. $x$ : int;
- Negative sign uses ~
- Operators: + - * div mod
- real
e.g. x: real;
- 3.45 or using e notation (3E7)
- Operators: + - */
- Conversion functions: real(integer), floor(real), abs(x)
- string e.g. s: string;
- Delimited by double quotes
- Caret $\wedge$ is concatenation e.g. "house" $\wedge$ "cat"
- Function size returns length of string
- Special characters: $\backslash n \backslash t \backslash " \backslash \backslash$
- bool
e.g. b: bool;
- true (and false)


## ML: operators

- All operators are infix
- Logical operators:
- Short-circuit evaluation
- if ... then ... else is an expression, not a control structure...

| not | Negation |
| :--- | :--- |
| andalso | Conjunction |
| orelse | Disjunction |
| if ... then ... else | conditional selection |

- Numeric operators:
- The usual $<,>,<=,>=$ and $<>$ are available
- For reals, = and $<>$ are not available ( $\mathrm{a}<=\mathrm{b}$ andalso $\mathrm{a}>=\mathrm{c}$ )
- For strings, these can be used for lexicographic ordering
- Operator overloading:
- Same symbol could be used for operations that are internally dissimilar
- *, +, - , <, <=, >, >= are all overloaded
- Leftmost argument is inspected first to decide on type


## ML: assignment

- Use val to assign value to variables
val <indentifier-name> = <expression>;
- Examples:
- val seconds $=60$;
$>$ val seconds $=60$ : int
- val minutes = 60;
$>$ val minutes $=60$ : int
- val tm $=$ seconds $*$ minutes;
$>$ val tm $=86400$ : int
- val shout = "aaa" ^ "rgh" ^ "!!!! ";
> val shout = "aaargh!!!" : string


## ML: constructor types - lists

- Syntax


## [ $\mathbf{o b j}_{1}, \mathbf{o b j}_{2}, \ldots$ ]

- Objects in a list must be homogenous (same type)
- E.g.

| $[1,2,3]$ | $:$ int list | -- |
| :--- | :--- | :--- |
| $[" d o g ", ~ " c a t ", ~ " m o o s e "] ~ l i s t ~ o f ~ i n t e g e r s ~$ |  |  |
| $[1.0,2.0,3.0]$ | : string list | $--a$ list of strings |
| $[[1,2],[3,4],[5,6]]$ | : real list | -- a list of reals |
|  | : int list list | $--a$ a list of lists of integers |

- The empty list is written [] or nil
- Operations:
- @ operator is used to concatenate two lists of the same type
- :: operator returns a new list with the first argument append to the front
- E.g. $2::[3,4]$ returns $[2,3,4]$
[1,2]::[ [3,4], [5,6]] returns ??
- hd returns the first element of a list
- E.g. hd[1,2,3] returns 1
- tl returns the tail
- E.g tl[1,2,3] returns [2,3]


## ML: constructor types - tuples

- Syntax


## ( $\mathbf{o b j}_{1}, \mathbf{o b j}_{2}, \ldots$ )

- Objects in a tuple can be heterogeneous (different types)
- E.g.

| (2, "abc") | : int * string |
| :--- | :--- |
| (2,3.0, "abc") | : int * real * string |
| $(2,(3.0$, "ab"), "cd") | : int * (real * string ) string |
| $[(1$, "a"),(3, "bc"),(7, "efg")] | : (int * string) list |

- The empty tuple is written () and often called unit
- Composite format of a tuple can be used on left-hand side of val
- Eg. - val (day, month, year ) = (13,"March",1066);
- Operations:
- \# operator to extract the ith field of a tuple
- \#2(6,7,"abc") returns 7
- \#3(6,7,"abc") returns abc
- = and <> operators for equality/in-equality

$$
\begin{aligned}
& - \text { val x = ~3; } \quad / /-3 \\
& -(3, " a ", \text { true })=(\text { abs x, "a", }(3>2)) ;
\end{aligned}
$$

## ML: constructor types - functions

- Syntax fun <func-name> <input-param> = <expression>;
- Keyword fun starts the function declaration
- Function arguments don't always need parentheses, doesn't hurt to use them
- Examples:
- fun fahrToCelsius $\mathrm{f}=$

$$
\text { (f -freezingFahr) * } 5 \operatorname{div} 9
$$

- fun foo $L=$

$$
(1+\mathrm{hd} \mathrm{~L})::(\mathrm{tl} \mathrm{~L}) ;
$$

## ML: constructor types - functions

- Functions have types too. ML interpreter will infer the type.
- E.g. fun square $x=$

$$
\text { x * } x
$$

- The function square takes an integer as input and returns an integer as output.
- This is written as square: int $\rightarrow$ int $(\rightarrow$ indicates this a function)


## ML: constructor types - functions

- ML figures out the input and/or output types for simple expressions, constant declarations, and function declarations
- Type checking requires that type expression of functions and their arguments match, and that type expression of context match result of function
- If the default isn't what you want, you can specify the input and output types

$$
\begin{array}{r}
\text { fun divideBy2 (y : real) }= \\
\text { y / 2.0; }
\end{array}
$$

- What is this doing?
fun foo (m, n) =
if $\mathrm{m}>\mathrm{n}$ then
[]
else $m$ :: foo( $\mathrm{m}+1, \mathrm{n}$ );
$>$ foo(1,6);


## ML: constructor types - functions

- Examples:
- Factorial (n!)

```
fun fact n =
if n = 0 then
    1
else
    n * fact(n-1);
```

```
fun 4 > 4* (fact 3)
    ->4*3*(fact 2)
    \rightarrow 4 * 3 * 2 ( f a c t ~ 1 )
    ->4*3*2*1*(fact 0)
    ->4*3*2*1*1
    ->24
```

- List reverse
fun reverse $\mathrm{L}=$
if $\mathrm{L}=$ nil then
nil
else
$\quad$ reverse(tl L) @ [hd L];


## ML: local environment using let

- Syntax

```
let
    val <variable > = <expression }\mp@subsup{}{1}{}>\mathrm{ ;
    val <variable e > = <expression n
in
    <expression>
end;
```

- Let allows declarations to be used in expressions
- Similar to Let* in Scheme
- Example:
- Compute hundredth power of a number

```
fun hundredthPower(x : real ) =
    let
        val four \(=\mathrm{x} * \mathrm{X} * \mathrm{x} * \mathrm{x}\);
    val twenty \(=\) four \(*\) four \(*\) four \(*\) four \(*\) four;
    in
    twenty* twenty* twenty* twenty* twenty
    end;
```


## ML: pattern matching

- Syntax

$$
\begin{aligned}
& \text { fun <func> <pattern }{ }_{1}>=\text { <expression }{ }_{1}> \\
& \text { | <func> <pattern }{ }_{2}>=\text { <expression }{ }_{2}> \\
& \text { <func> <pattern }{ }_{\mathrm{n}}>=\text { <expression }{ }_{\mathrm{n}} \text { > }
\end{aligned}
$$

- Define a function by a series of equations, LHS is a pattern.
- Always put the most specific pattern first
- ML interpreter will use the first equation whose LHS matches
- Example:
- Fibonacci function $\left(a_{n}=a_{n-1}+a_{n-2}:-\quad 0,1,1,2,3,5,8,13,21, \ldots\right)$
fun fib $n=$

$$
\begin{aligned}
& \text { if } \mathrm{n}=0 \text { then } 0 \\
& \text { else if } \mathrm{n}=1 \text { then } 1 \\
& \quad \text { else fib(n-1) }+\mathrm{fib}(\mathrm{n}-2) \text {; }
\end{aligned}
$$

$$
\begin{aligned}
& \text { fun } \mathrm{fib}(0)=0 \\
& \mid \quad \operatorname{fib}(1)=1 \\
& \quad \mathrm{fib}(\mathrm{~N})=\mathrm{fib}(\mathrm{~N}-1)+\mathrm{fib}(\mathrm{~N}-2) ;
\end{aligned}
$$

- Pattern matching is powerful:
- Allows the programmer to see the arguments. No more hd's and tl's.


## ML: pattern matching - cont'd

- Examples:
- Sum all the elements in a list of integers fun listsum $\mathrm{L}=$
if (null L) then 0 else (hd L) + listsum(tl L);

Better
fun listsum [] $=0$
| listsum L = (hd L) + listsum(tl L);
listsum[1,2,3,4] $\rightarrow 1+$ listsum[2,3,4]
$1+2+$ listsum $[3,4]$
$1+2+3+$ listsum[4]
1+2+3+4+listsum[]
$1+2+3+4+0$
10

## ML: pattern matching - cont'd

- Examples:
- Reversing a list fun reverse $\mathrm{L}=$
if $\mathrm{L}=$ nil then nil
else reverse(tl L) @ [hd L];
fun reverse(nil) $=$ nil
| reverse(h::t) = reverse(t) @ [h];
reverse $[1,2,3] \rightarrow$ reverse[2,3] @ 1
$\rightarrow$ reverse[3] @ 2 @ 1
$\rightarrow$ reverse[]@3@2@1
$\rightarrow[]$ @ 3 @ 2 @ 1
$\rightarrow[3,2,1]$


## ML: pattern matching - cont'd

- Examples:
- Return first $n$ elements of a list
fun take ([ ], Index) = [ ]
| take (h::tl, Index) = if Index > 0 then
h::take(tl, Index - 1) else
[ ];

$$
\begin{aligned}
\operatorname{take}([1,2,3], 2) & \rightarrow 1:: \text { take }([2,3], 1) \\
& \rightarrow 1:: 2 \text { take([3], } 0) \\
& \rightarrow 1:: 2::[] \\
& \rightarrow[1,2]
\end{aligned}
$$

