



Principles of Programming Languages

Lecture 13

Wael Aboulsaadat

wael@cs.toronto.edu

<http://portal.utoronto.ca/>

Acknowledgment: parts of these slides are based on material by Diane Horton & Eric Joanis @ UoT

References: Scheme by Dybvig

PL Concepts and Constructs by Sethi

Concepts of PL by Sebesta

ML for the Working Prog. By Paulson

Prog. in Prolog by Clocksin and Mellish

PL Pragmatics by Scott

ML: pattern matching – cont'd

- **Patterns may consist of constants (integers , true, false..) , tuples and variables. Arithmetic or logical expressions are invalid.**
 - E.g. `fun wrong(x=y) = “...”`
- **No duplicates in patterns**
 - E.g. `fun wrong_equal (x,y) = true`
`| wrong_equal (x,y) = false;`
- **Pattern matching with wild cards**
 - E.g. `fun first (x,_) = x;`
 - Matches anything like a variable. Binds nothing.
 - Avoid need to name every pattern
- **ML does extensive pattern checking**
 - E.g. `- fun reverse (h::t) = reverse(t) @ [h];`
`> Warning: match nonexhaustive`

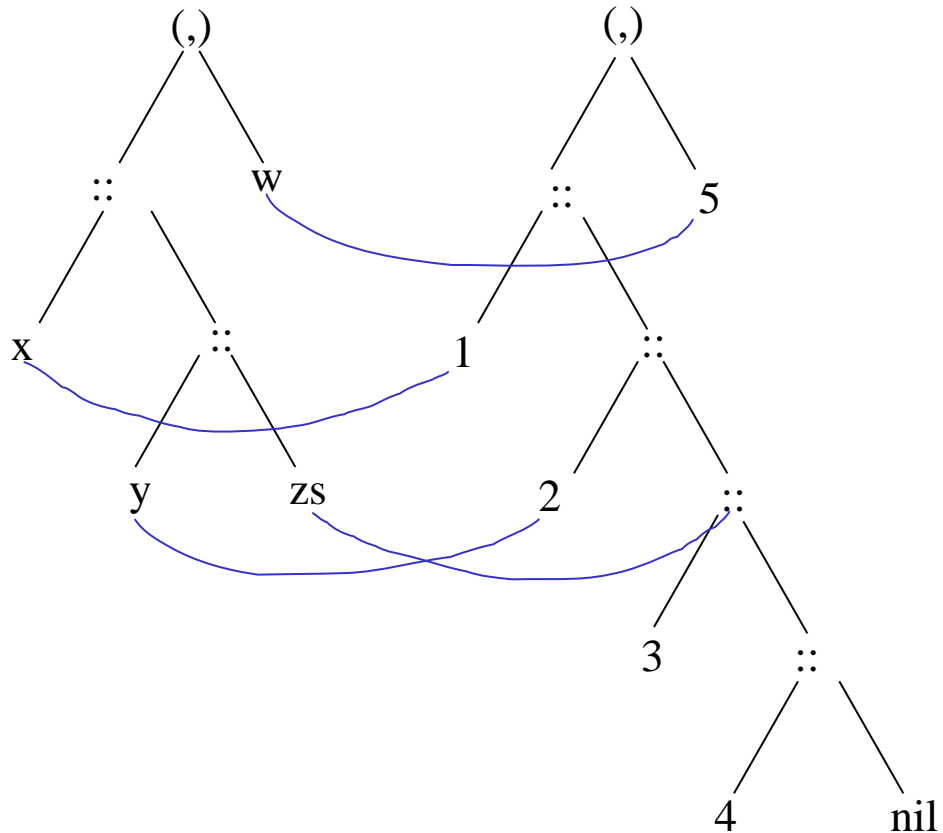
ML: pattern matching – cont'd

- How ML matches patterns?

```

fun func (x::y::zs,w) =
  | (...) =
  | (...) =
  
```

```
func([1,2,3,4],5)
```



ML: record type

- We have seen list [] and tuple ()...

- Record syntax

$\{ \langle \text{field}_1 \rangle : \langle \text{type}_1 \rangle, \langle \text{field}_2 \rangle : \langle \text{type}_2 \rangle, \dots, \langle \text{field}_n \rangle : \langle \text{type}_n \rangle \}$

- A record instance is defined as

$\{ \langle \text{label}_1 \rangle = \langle \text{value}_1 \rangle, \langle \text{label}_2 \rangle = \langle \text{value}_2 \rangle, \dots, \langle \text{label}_n \rangle = \langle \text{value}_n \rangle \}$

- A record is a structured data type in which each element is accessed by a unique name.

– E.g. { name: string,
age: int,
salary: int }

– { name = “Dave”, age = 77, salary=99000 }

> val it = { name = “Dave”, age = 77, salary=99000 } : { name:string,
age:int,salary:int }

A field called name with type string
A field called age with type int
A field called salary with type int



ML: record type – cont'd

- **Operations**

- # operator to extract a field from a record instance

- E.g.

- #salary {name =“john”, age=35, salary=90};

- > val it = 90 : int

- #options {startcity="toronto",endcity="boston",
options=("12",10,"K")};

- > val it = ("12",10,"K") : string * int * string



ML: record type – cont'd

- **Named Types**

- ML provides a way to give a name to a type

- E.g. - `type waitress = { name: string , wages: int, tips: int };`

- Named types can be used anywhere that ML types can be.

- E.g. - `fun income (w: waitress) =`

- `#wages w + #tips w;`

- `> val income = fn : waitress -> int`

- `fun income (w: { name: string , wages: int, tips: int }) =`

- `#wages w + #tips w;`

- `> val income = fn : { name:string, tips:int, wages:int} -> int`

- Named types can be used in type declaration

- E.g. - `type waitresses = waitress list;`

- `[{name="sally", wages= 20, tips=10},
{name="alice", wages= 15, tips=15}]`



ML: record type – cont'd

- **Named Types – cont'd:**

- E.g.: *finding the total income of all waitresses*

- > `type waitress = { name: string , wages: int, tips: int };`

- > `fun income (w: waitress) =`

- `#wages w + #tips w;`

- > `type waitresses = waitress list;`

- > `fun total (WL: waitresses) =`

- `if WL = [] then 0`

- `else income(hd WL) + total(tl WL);`

- WL = [{name="sally", wages= 20, tips=10},

- {name="alice", wages= 15, tips=20},

- {name="sue", wages= 25, tips=20}]

30 +

[{name="alice", wages= 15, tips=20}

{name="sue", wages= 25, tips=20}]

30 + 35 +

[{name="sue", wages= 25, tips=20}]

30 + 35 + 45

-total(WL)



110

ML: pattern matching on records

- **Recall syntax**
`fun <func> <pattern1> = <expression1>`
`| <func> <pattern2> = <expression2>`
`.....`
`| <func> <patternn> = <expressionn>`
- **You can use patterns to match on records**
 - E.g. *finding the total income of all waitresses*
`> type waitress = { name: string , wages: int, tips: int };`
`> fun income (w: waitress) =`
`#wages w + #tips w;`
`> type waitresses = waitress list;`
`> fun total ([]: waitresses) = 0`
`| total(W::WLTail) = (income W) + (total WLTail);`
- **You can also use wild cards ...**
 - E.g.
`> fun costly({price:int, ...}: footype) = price > 100.0;`

ML: ...

- What does ML infer about this function?

```
- fun length L =  
    if (null L) then 0  
    else 1 + length(tl L);
```

```
- length[1,2,3,4]
```

```
> val it = 4 : int
```

```
- length["ab","cd","xy"];
```

```
> val it = 3 : int
```

```
- length[[1,2],[3,4],[123,123,222],[1]];
```

```
> val it = 4 : int
```

- Seems length has/accept these types

- int list → int
- String list → int
- int list list → int

- Obviously, we would like *length* to apply to any kind of list.

ML: ...

- What does ML infer about this function?

- **fun** length L =
 if (null L) **then** 0
 else 1 + length(**tl** L);

- In ML, *length* has all of these types. This is written as

- length: 'a list → int**

- 'a is a type variable. It stands for any type
 - This means that the input to length is a list of items all of type 'a where 'a can be **int**, **string**, **int list**, or any other type.

- In fact, that's what ML infers for this function

- **fun** length L = **if** (null L) **then** 0 **else** 1 + length(**tl** L);
> val length = fn : 'a list -> int



ML: polymorphism

- **Greek:** *poly* = many , *morph* = form
- **Definitions:**
 - **Polymorphism:**
 - dictionary.com: the capability of assuming different forms; the capability of widely varying in form. The occurrence of different forms, stages, or types
 - Software: a value/variable can belong to multiple types
 - **Monomorphism:**
 - Dictionary.com: having only one form, same genotype...
 - Software: every value/variable belongs to exactly one type
- **Why is useful?**
 - To avoid redundant function definitions, e.g.:
 - int-length : `int list` → `int`
 - real-length: `real list` → `int`
 - string-length: `string list` → `int`
 - code for each of these functions would be virtually identical!
 - Polymorphism adds flexibility & great convenience.... but...

ML: polymorphism types

- **Ad-hoc polymorphism:**

- Different operations on different types known by the same name (*also called overloading*)
- E.g. $3 + 4$ vs. $3.1 + 4$ *compiler/interpreter must change 4 to 4.0 first*
vs. "ab" + "cd"

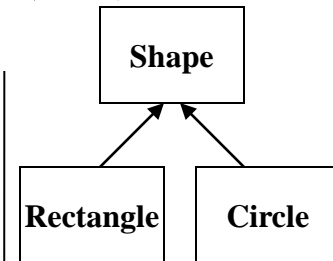
- **Inheritance polymorphism:**

- Use sub-classing to define new versions of existing functions (*OO*)

- E.g.:

```
public class Shape{  
  
    public void draw( int x, int y){  
        // do nothing  
    }  
}  
  
public class Rectangle extends Shape{  
    public void draw( int x, int y){  
        // draws a rectangle  
    }  
}
```

```
public class Circle extends Shape{  
    public void draw( int x, int y){  
    }  
}  
....  
Shape myShape;  
myShape = new Rectangle( );  
  
..... // some other piece of code  
myShape.draw( );
```





ML: polymorphism types – cont'd

- **Parametric Polymorphism(ML):**
 - Allows types to be parameters to functions and other types.
 - Basic idea is to have a type variable...
 - Type of function depend on type of parameter
 - Implementation (ML):
 - One copy of code is generated
 - Polymorphic parameters must internally be implemented as pointers

ML: polymorphism – cont'd

- Polymorphic functions are very common in ML:

- fun id X = X;
> val id = fn : 'a -> 'a

- id 7;
> val it = 7 : int
- id "abc";
> val it = "abc" : string

- fun listify X = [X];
> val listify = fn : 'a -> 'a list

- listify 3;
> val it = [3] : int list
- listify 7.3;
> val it = [7.3] : real list

- fun double X = (X,X);
> val double = fn : 'a -> 'a * 'a

- double "xy";
> val it = ("xy","xy") : string * string
- double [1,2,3];
> val it = ([1,2,3],[1,2,3]) : int list * int list

- fun inc(N,X) = (N+1,X);
> val inc = fn : int * 'a -> int * 'a

- inc (4,(34,5));
val it = (5,(34,5)) : int * (int * int)

ML: polymorphism – cont'd

- **Polymorphic functions are very common in ML:**

<pre>- fun swap(X,Y) = (Y,X); > val swap = fn : 'a * 'b -> 'b * 'a</pre> <hr/>	<pre>- swap ("abc",7); > val it = (7,"abc") : int * string - swap (13.4,[12,3,3]); val it = ([12,3,3],13.4) : int list * real</pre> <hr/>
<pre>- fun pair2list(X,Y) = [X,Y]; > val pair2list = fn : 'a * 'a -> 'a list</pre> <hr/>	<pre>- pair2list(1,2); > val it = [1,2] : int list - pair2list(1,"cd"); ?</pre> <hr/>
<pre>- fun apply(Func,X) = Func X; > val apply = fn : ('a -> 'b) * 'a -> 'b</pre> <hr/>	<pre>- apply (hd, [1,2,3]); > val it = 1 : int - apply (length, [23,100]); > val it = 2 : int</pre> <hr/>
<pre>- fun applytwice(Func,X) = Func(Func X); > val applytwice = fn : ('a -> 'a) * 'a -> 'a</pre>	<pre>- applytwice (square,3); > val it = 81 : int - applytwice (tl, [1,2,3,4]); - ?</pre>

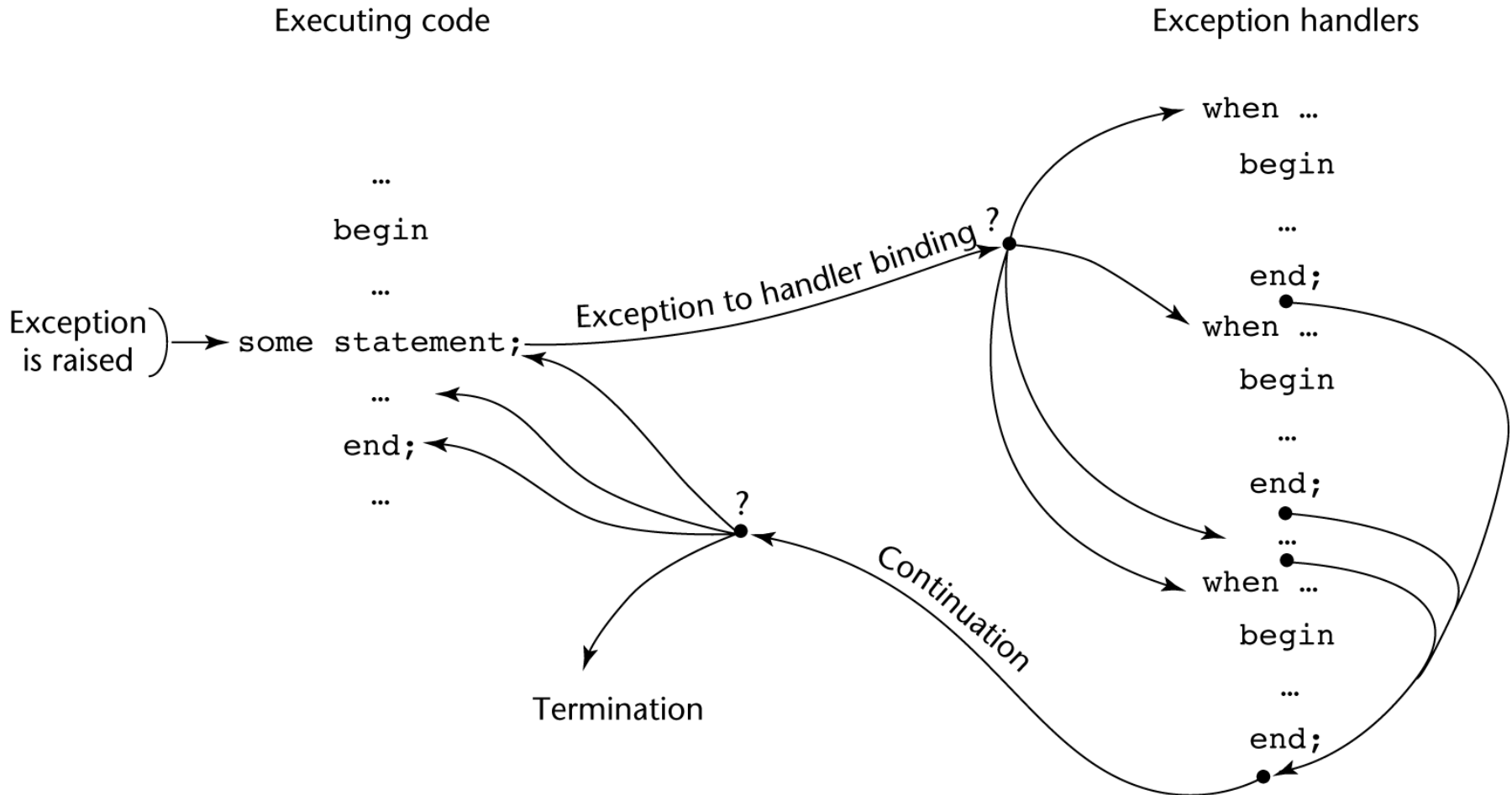
ML: polymorphism – cont'd

- **Operators that restrict polymorphism**
 - Arithmetic operators: `+`, `-`, `*` and `-`
 - Division-related operations such as `/`, `div` and `mod`
 - Inequality comparison operators: `<`, `<=`, `>=`, and `>`
 - Boolean connectives: `andalso`, `orelse` and `not`
 - String concatenation operator: `^`
 - Type conversion operators
 - E.g. `ord`, `chr`, `real`, `str`, `floor`, `ceiling`, `round`, `truncate`,...
- **Operators that allow polymorphism**
 - Tuple operators
 - List operators
 - Equality operators `=` and `<>`

Exceptions: introduction

- An **exception** is any unusual event, erroneous or not, that is detectable either by hardware or software and that may require special processing.
- The special processing that may be required by the detection of an exception is called **exception handling**. This processing is done by a code unit called the **exception handler**.
- **Why do we need exceptions if the language is strongly typed?**
 - In a language without exception handling: when an exception occurs, control goes to the operating system, where a message is displayed and the program is terminated.
 - In a with exception handling: programs are allowed to trap some exceptions, thereby providing the possibility of fixing the problem and continuing.

Exceptions: execution flow



Exceptions: why?

- **How was error handling done in early programming languages?**
 - Send an auxiliary parameter or use the return value to indicate the return status of a subprogram (e.g. C standard library functions)

```
nError = mult(matrix1,matrix2,product);  
if( nError == -1){ // error  
  
}  
else{ // no error, continue normally  
  
}
```

- Pass a label parameter to the subprogram. If an error occurs, use the label to jump to another location in the program (e.g. FORTRAN)

```
mult(matrix1,matrix2,product,label)  
  
...  
if error  
goto label
```

- Pass an error-handler subprogram to the called subprogram.

```
mult(matrix1,matrix2,product,error_func)  
  
if error  
error_func(...)
```

ML: exceptions

- **Syntax** `exception <exception-name> of <type-expression>`

- **Example:**
 - `exception NegArg of int;`
> `exception NegArg of int`

 - `fun fact N = if N = 0 then 1`
 `else if N > 0 then N * fact(N-1)`
 `else raise NegArg(N);`
> `val fact = fn : int -> int`

 - `fact(5);`
> `val it = 120 : int`

 - `fact(~5);`
> `uncaught exception NegArg raised at: ...`