

Principles of Programming Languages Lecture 18

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Function Type

• Function types

- A type that allow an object to be invoked or called as if it were an ordinary function
- E.g. Sorting an array containing a user defined type

```
// Declaration of C sorting function
void sort (int [] itemlist, int numitems, int (*cmpfunc)(item*, item*) );
...
```

```
// Callback function
int compare_function( item* A, item* B)
{
    //.....
//.....
}
```

sort(itemlist, numitems, compare_function);



Function Type

• Function types

- A type that allow an object to be invoked or called as if it were an ordinary function
- E.g. using different sorting algorithm

// Declaration of C sorting function

void bubblesort (int [] itemlist, int numitems, int (*cmpfunc)(int*, int*)) {

```
}
void quicksort (int [] itemlist, int numitems, int (*cmpfunc)(int*, int*) ) {
....
void doSomething(....) {
    ...
    func1(itemlist , numitems , bubblesort);
    ...
    func2(itemlist , numitems , quicksort);
}
void doSomething(....) {
    ...
    func2(itemlist , numitems , quicksort);
}
void func1(....,sortfunc ) {
    sortfunc(itemlist,numitems)
```



Data types Summary

- Primitive types
 - Integer, Float, Boolean, Char
- Structured Types
 - Strings, Ordinal, Arrays, Associative Arrays, Records, Union, Lists
- Object Type
- Class Type
- Function Type

Note: Don't confuse data types with API data structures (BST, Graphs,...)

Type Conversions

• Often want to write expressions that are mixed mode (contain operands of more than one type)

```
real + integer
```

- This implies a need to convert one type to another so the expression can be evaluated
 - coercion is *implicit conversion*, done automatically by the compiler
 - implies semantics that define rules for determining type to convert to from operands
 - problem is loss of error detection
 - casts are <u>explicit conversions</u> specified by the programmer
 - can lead to very clumsy expressions if doing a lot of mixed mode expressions

int (Index);	// Ada and Pythor
(int) Index;	// C and Java



- Whether implicit or explicit, conversions can be
 - widening
 - convert to a type with a greater representation range
 - although perhaps with a loss of precision
 - integer \rightarrow real
 - narrowing
 - convert to a type with a more restricted range
 - double precision \rightarrow real
 - promoting
 - convert to a type with additional semantic information
 - integer \rightarrow character e.g. NewChar = chr(IntValue)
 - demoting
 - strip away semantic information
 - character \rightarrow integer



- PL/I allows coercion between almost any types
 - DCL A, B, C INT;
 - if $(A \leq B \leq C)$ then ...
 - A <= B yields a Boolean value, which is a single bit 0 or 1
 - convert bit to integer to compare to C
 - bit <= C is true for any positive values of C
- Ada allows no coercion
 - all conversions must be casts
 - conversions are allowed between all numeric types
 - other conversions only allowed between derived types that share an ancestor

type foo is new Boolean; type bar is new Boolean; A: foo; B: bar; A := foo(B);



Data Types Questions

- What are the data types in the language?
- How is the data type declared?
- What operations are allowed on each data type?
- How to reference the data type?
- How is the data type implemented by the compiler/interpreter?
- What conversion rules exist for each datatype?
- What type checking does the language support?

Data types: how many should a language have?



• Early programming languages:

- Many data types
- Support large range of applications

• Modern programming languages:

- Few basic types and few basic data structure
- Allow a programmer to design complex structures for every need

Components of an Imperative Language

- Data types
- → Variables & Expressions
- Assignment construct
- Iteration construct
- Branching construct
- Function construct
- Container construct



Names: design decisions

- How long can a name be? What characters can be used? Are connecters allowed?
 - 30-60 max characters are most practical and used by most languages
- Are names case sensitive?
 - Disadvantage: names that look alike are different
 - Java and Modula-2 are the worst because predefined function names are mixed (
 E.g. java.lang.Integer.parseInt(...)) You just have to remember that...
- What is the set of keyword vs. reserved words?
 - Keywords: special only in certain contexts
 - E.g. In Fortran

REAL TEMP	// ok
REAL = 3.4	// also ok!

- *Reserved words*: cannot be used by programmer as names

• E.g. In Java double int = 5; // error



Variables: introduction

• A variable is an abstraction of a computer memory cell

- A variable is a not a name!
- A variable can be characterized by a sixtuple of attributes:
 - Name
 - Address:
 - A variable may have different addresses at different times during execution
 - A variable may have different addresses at different places in a program
 - If two variable names can be used to <u>access the same memory location</u>, they are called *aliases (aliases are harmful to readability)*
 - Value:
 - The content of the location with which the variable is associated.
 - Type:
 - Determines the range of values of variables and the set of operations that are defined for values of that type; in the case of floating point, type also determine the precision.
 - Lifetime:
 - It is the time during which the variable is bound to a particular memory cell
 - Scope



Variables: value vs. reference

- How will the variables be used in statements and how they are represented in memory?
- Example:
 - // Java language
 int b = 3;
 int a = b;

// C language int c = 3; int b = &c; int a = b; int* z; z = (int)malloc(size(int)) *z = 10



Variables: value vs. reference

• Example:



Variables: variations

• Value-oriented

- C/C++, Pascal, Basic, Ada

Reference-oriented

- C/C++, Scheme/ML (functional languages)
- Clu, Smalltalk
- Hybrids
 - Algol-68, C/C++
 - Java
 - built-in types are values (int, float, double..)
 - user-defined types are objects (i.e., references)



Variables: variations

- Encoded variables
- E.g. Perl
 - scalars starts with \$

\$numberOfRooms = 23;

– Arrays starts with @

@stringArray = ("This", "is", 'an', "array", 'of', "strings");

Associative Array starts with %

%associativeArray = ("Jack A.", "Dec 2", "Joe B.", "June 2", "Jane C.", "Feb 13");



Binding: introduction

- A binding is an association between two things, such as a name and the thing it names.
 - E.g.
 - int x; // for this to work the compiler must bind a memory // cell (that is sufficient to hold an integer) with the // identifier x so later on the programmer can write

x = 5;

- The *binding time* is the point at which a binding is created or, more generally, the point at which any implementation decision is made.
 - When will the name bind to the value?



Binding: variables to storage

• Binding creation/allocation:

- Getting a (memory) cell from some pool of available (memory) cells and establishing an association between it a specific variable

• Binding destruction/de-allocation:

 Putting a (memory) cell back into the pool and destroying the association between it and the variable.

• Binding lifetime:

– The period of time from creation of a binding to its destruction



Binding: when?

- During language design
 - E.g.: In Pascal, % operator is bound to mod operation

• During language implementation

- E.g. In Java, data type double is bound to certain range

• At compile time

– E.g. In Java, int x; x is bound to particular data type (integer)

• When linking

- E.g. A call to subprogram *foo* in a separate library (.dll / .lib / .so) is bound to the subprogram code

• At load time

– E.g. In C, a variable may be bound to a storage cell when the program is loaded

• At run time

 E.g. In C++, a variable may be bound to a storage cell after the program is loaded (using pointers).



Binding: classification

- Static vs. Dynamic:
 - The terms *static* and *dynamic* are generally used to refer to things bound before run time and at run time, respectively.
 - Static binding:
 - Can mean many different times (e.g., language design, compile time, etc.).
 - Dynamic binding:
 - Generally referring to binding times such as when variable values are bound to variables.
 - Advantage: flexibility
 - Disadvantage: High cost, Type error detection by the compiler is difficult





- It is the activity of ensuring that the operands of an operator are of • compatible types:
 - Subprograms are considered as operators
 - A compatible type is one that is either legal for the operator or is allowed under language rules to be implicitly converted to a legal type:
 - E.g.

float R = (float)10; // valid in most languages int X = (int)10.5; // invalid in most languages

Type error: ۲

- It is an attempt to apply a function to an argument of the wrong type

– E.g.

int X = java.lang.StrictMath.round("nancy");





- Type checking depends on binding:
 - *Static type checking:* if all bindings of variables are done at compile type, then type checking can be done statically.
 - Dynamic type checking: if some bindings of variables are done at run time, then type checking will be for thoese variables dynamically, i.e. when the program is running.



Type Checking: classification

• Strongly typed language:

- It is one in which each name in a program in the language has a <u>single</u> type associated with it, and that type is known at compile time (i.e. statically bound).
- E.g.: ML, Pascal, Ada

• Not-Strongly typed language:

- It is one in which variable types may be known but the storage location to which it is bound may store values of different types at different times.
- E.g.: C, C++, Fortran

• Weakly typed language:

- It is one in which a name in a program in the language can change the type associated with it during run time and type checking is
- E.g.: Basic, Perl, Python

Components of an Imperative Language

- Data types
- Variables
- → Operators & Expressions
- Iteration construct
- Branching construct
- Function construct
- Container construct



Operators & Expressions

- Mathematical operators
- Logical operators
- Bitwise operators
- User-defined operators



Mathematical & Logical Operators

- Precedence rules
 - Norm: respect mathematical precedence
 - Force evaluation of specific terms using ()

$$-$$
 E.g. $(x + y) * z$

 $\frac{vs.}{x + (y * z)}$

• Short circuit evaluation

Bitwise Operators



- Low level operations supported directly by CPU
- Performed on register content
- Bitwise AND, Bitwise OR, Bitwise, Shift left, Shift right, XOR,...

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Bitwise Operators

- Why care?
- Languages: C/C++, PhP
- E.g.

```
// C language. left shifting is the equivalent of multiplying by
// a power of two
int mult_by_pow_2(int number, int power) {
    return number<<pre>power;
}
```