## Graphs



## Graphs

- A graph is a pair $(\boldsymbol{V}, \boldsymbol{E})$, where
- $V$ is a set of nodes, called vertices
- $E$ is a collection of pairs of vertices, called edges
- Vertices and edges are positions and store elements
- Example:
- A vertex represents an airport and stores the three-letter airport code
- An edge represents a flight route between two airports and stores the mileage of the route



## Edge Types

- Directed edge
- ordered pair of vertices $(\boldsymbol{u}, \boldsymbol{v})$
- first vertex $\boldsymbol{u}$ is the origin
- second vertex $v$ is the destination

- e.g., a flight
- Undirected edge
- unordered pair of vertices (u,v)
- e.g., a flight route

- Directed graph
- all the edges are directed
- e.g., route network
- Undirected graph
- all the edges are undirected
- e.g., flight network


## Applications

- Electronic circuits
- Printed circuit board
- Integrated circuit
- Transportation networks
- Highway network
- Flight network
- Computer networks
- Local area network
- Internet
- Web
- Databases

- Entity-relationship diagram


## Terminology

- End vertices (or endpoints) of an edge
- U and V are the endpoints of a
- Edges incident on a vertex
- $a, d$, and b are incident on V
- Adjacent vertices
- U and V are adjacent
- Degree of a vertex
- X has degree 5
- Parallel edges
- $h$ and $i$ are parallel edges
- Self-loop

- j is a self-loop


## Terminology (cont.)

- Path
- sequence of alternating vertices and edges
- begins with a vertex
- ends with a vertex
- each edge is preceded and followed by its endpoints
- Simple path
- path such that all its vertices and edges are distinct
- Examples
- $P_{1}=(V, b, X, h, Z)$ is a simple path
- $P_{2}=(U, C, W, e, X, g, Y, f, W, d, V)$ is a
 path that is not simple


## Terminology (cont.)

- Cycle
- circular sequence of alternating vertices and edges
- each edge is preceded and followed by its endpoints
- Simple cycle
- cycle such that all its vertices and edges are distinct
- Examples
- $C_{1}=(V, b, X, g, Y, f, W, c, U, a, \cdot)$ is a simple cycle
- $\mathrm{C}_{2}=(\mathrm{U}, \mathrm{c}, \mathrm{W}, \mathrm{e}, \mathrm{X}, \mathrm{g}, \mathrm{Y}, \mathrm{f}, \mathrm{W}, \mathrm{d}, \mathrm{V}, \mathrm{a},-\mathrm{J})$ is a cycle that is not simple



## Properties

Property 1
$\Sigma_{v} \operatorname{deg}(\boldsymbol{v})=2 m$
Proof: each edge is counted twice

## Notation

| $\boldsymbol{n}$ | number of vertices |
| :---: | :--- |
| $\boldsymbol{m}$ | number of edges |
| $\operatorname{deg}(\boldsymbol{v})$ | degree of vertex $v$ |

Property 2
In an undirected graph with no self-loops and no multiple edges

$$
m \leq n(n-1) / 2
$$

Proof: each vertex has degree at most $(\boldsymbol{n}-1)$

What is the bound for a


## Example

- $n=4$
- $m=6$
- $\operatorname{deg}(\boldsymbol{v})=3$


## Main Methods of the Graph ADT

- Vertices and edges
- are positions
- store elements
- Accessor methods
- endVertices(e): an array of the two endvertices of e
- opposite( $\mathrm{v}, \mathrm{e}$ ): the vertex opposite of v on e
- areAdjacent $(\mathrm{v}, \mathrm{w})$ : true iff v and $w$ are adjacent
- replace $(\mathrm{v}, \mathrm{x})$ : replace element at vertex $v$ with $x$
- replace(e, x): replace element at edge $e$ with $x$
- Update methods
- insertVertex(0): insert a vertex storing element o
- insertEdge( $\mathrm{v}, \mathrm{w}, \mathrm{o}$ ): insert an edge ( $\mathrm{v}, \mathrm{w}$ ) storing element o
- removeVertex(v): remove vertex $v$ (and its incident edges)
- removeEdge(e): remove edge e
- Iterable collection methods
- incidentEdges(v): edges incident to v
- vertices(): all vertices in the graph
- edges(): all edges in the graph


## Edge List Structure

- Vertex object
- element
- reference to position in vertex sequence

- Edge object
- element
- origin vertex object
- destination vertex object
- reference to position in edge sequence
- Vertex sequence
- sequence of vertex objects
- Edge sequence
- sequence of edge objects



## Adjacency List Structure

- Edge list structure
- Incidence sequence for each vertex

- sequence of references to edge objects of incident edges
- Augmented edge objects
- references to associated positions in incidence sequences of end vertices



## Adjacency Matrix Structure

- Edge list structure
- Augmented vertex objects

- Integer key (index) associated with vertex
- 2D-array adjacency array
- Reference to edge object for adjacent vertices
- Null for non nonadjacent vertices
- The "old fashioned" version just has 0 for no edge and 1 for edge



## Performance

| $\boldsymbol{n} \boldsymbol{n}$ vertices, $\boldsymbol{m}$ edges <br> - no parallel edges <br> - no self-loops | Edge <br> List | Adjacency <br> List | Adjacency <br> Matrix |
| :--- | :---: | :---: | :---: |
| Space | $\boldsymbol{n + \boldsymbol { m }}$ | $\boldsymbol{n + \boldsymbol { m }}$ | $\boldsymbol{n}^{2}$ |
| incidentEdges $(\boldsymbol{v})$ | $\boldsymbol{m}$ | $\operatorname{deg}(\boldsymbol{v})$ | $\boldsymbol{n}$ |
| areAdjacent $(\boldsymbol{v}, \boldsymbol{w})$ | $\boldsymbol{m}$ | $\min (\operatorname{deg}(\boldsymbol{v}), \operatorname{deg}(\boldsymbol{w}))$ | 1 |
| insertVertex $(\boldsymbol{o})$ | 1 | 1 | $\boldsymbol{n}^{2}$ |
| insertEdge $(\boldsymbol{v}, \boldsymbol{w}, \boldsymbol{o})$ | 1 | 1 | 1 |
| removeVertex $(\boldsymbol{v})$ | $\boldsymbol{m}$ | $\operatorname{deg}(\boldsymbol{v})$ | $\boldsymbol{n}^{2}$ |
| removeEdge $(\boldsymbol{e})$ | 1 | 1 | 1 |

